

Basis Theorems and Models of WKL_0

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Definition

\mathcal{L}_2 : The language of 2nd order arithmetic.

$M = (\mathbb{N}^M, S)$: \mathcal{L}_2 -structure. \mathbb{N}^M : the 1st order part of M ,
 S : the 2nd order part of M .

Note: We use ω to denote the standard natural numbers, and \mathbb{N} to denote the 1st part of a \mathcal{L}_2 -structure. If $\mathbb{N}^M \simeq \omega$ then M is called ω -model.

Subsystems of 2nd order arithmetic

$\text{RCA}_0 = \text{I}\Sigma_1^0 + \text{recursive sets exist.}$

$\text{WKL}_0 = \text{RCA}_0 + \text{any infinite binary tree has a path.}$

$\text{ACA}_0 = \text{RCA}_0 + \text{arithmetically definable sets exist.}$

General Form of Basis Theorems

Basis theorems have the following form in general:

Let T be a computable infinite binary tree. Then T has a path s.t. (some conditions).

Examples of Basis Theorems:

Low Basis Theorem(LBT)

T has a low path.

Hyperimmune-free Basis Theorem(HFBT)

T has a hyperimmune-free path.

Fact

There is a computable tree T s.t. each path of T is regarded as a countable ω -model of WKL_0 .

Combining this fact and basis theorems, we can make a model of WKL_0 with some properties:

Fact

There is an ω -model of WKL_0 which includes only low/hyperimmune-free sets.

In this talk,

- (1) we formalize this argument in 2nd order arithmetic,
- (2) decide the strength of low/hyperimmune-free basis theorem.

Definition(low set)

For $A \subseteq \omega$, A' is one of Σ_1^0 -complete set.

We say that a set A is low if $A' \leq_T \emptyset'$.

The jump operator is formalized in 2nd order arithmetic as follows: Let $\Phi(e, m, A)$ be a Σ_1^0 -universal formula. For $A \subseteq \mathbb{N}$, $A' = \{(e, m) : \Phi(e, m, A)\}$.

Definition(hyperimmune-free set)

Let $X \subseteq \omega$. X is hyperimmune-free if

$\forall f \leq_T X \exists g \leq_T \emptyset (f < g)$.

From now on, we assume T to be a binary tree.

Relativized Low Basis Theorem

ACA_0 proves

$$\forall X \forall T \leq_T X (|T| = \infty \rightarrow \exists Y \in [T] (Y \oplus X)' \leq_T X').$$

Relativized Hyperimmune-free Basis Theorem

ACA_0 proves

$$\forall X \forall T \leq_T X (|T| = \infty \rightarrow \\ \exists Y \in [T] (\forall f \leq_T Y \exists g \leq_T X (f < g))).$$

So, they are strong versions of weak König's lemma.

Lemma 1

In WKL_0 , $\exists X \Pi_1^0$ is also Π_1^0 .

Intuitively, a Π_1^0 formula $\varphi(X)$ corresponds to a effectively closed set of a Cantor space. Thus we can find a binary tree T s.t. $[T] = \{X : \varphi(X)\}$. Moreover this argument can be formalized in RCA_0 .

[Proof] Let $\varphi(X)$ be a Π_1^0 formula and T be a binary tree s.t. $\forall X (X \in [T] \leftrightarrow \varphi(X))$. Then

$$\text{WKL}_0 \vdash \exists X \varphi(X) \leftrightarrow |T| = \infty.$$

Clearly $|T| = \infty$ is Π_1^0 . \square

Lemma 2

WKL₀ proves the compactness of the Cantor space $2^{\mathbb{N}}$.

That is, for any Π_1^0 formula $\varphi(X, n)$,

$$\text{WKL}_0 \vdash (\forall n \exists X \forall i < n \varphi(X, i)) \rightarrow \exists X \forall n \varphi(X, n).$$

Intuitively, each n defines a closed subset \mathcal{C}_n of $2^{\mathbb{N}}$ s.t.

$\mathcal{C}_n = \{X : \varphi(X, n)\}$. Hence, this theorem says that if

$\{\mathcal{C}_n\}_{n \in \mathbb{N}}$ has finite intersection property, then $\bigcap_n \mathcal{C}_n \neq \emptyset$.

This can be proved by similar way of Lemma 1.

Low basis theorem (ACA_0)
$$\forall X \forall T \leq_T X (|T| = \infty \rightarrow \exists Y \in [T] (Y \oplus X)' \leq_T X').$$

By universal Σ_1^0 formula, we can take an enumeration

$\{\varphi_e(m, A)\}_{e \in \mathbb{N}}$ of all Π_1^0 formulas. Then,

$\text{RCA}_0 \vdash \forall A (A' \equiv_T \{(e, m) : \varphi_e(m, A)\})$.

We will show that

$$\exists Y \in [T] (\{(e, m) : \varphi_e(m, Y \oplus X)\} \leq_T X').$$

Claim: $\exists Y \in [T](\{(e, m) : \varphi_e(m, Y \oplus X)\} \leq_T X')$.

[Proof] By using arithmetical comprehension, we can define a maximal subsequence $\{(e_i, m_i)\}_i$ of \mathbb{N}^2 s.t.

$$\begin{aligned} \exists i \forall Z (\varphi_{e_i}(n_i, Z \oplus X) \leftrightarrow Z \in [T]), \\ \forall k \exists Z \forall i < k (\varphi_{e_i}(m_i, Z \oplus X)). \end{aligned}$$

By lemma 1, $\{(e_i, m_i)\}_{i \in \mathbb{N}}$ is computable from X' .

Now there exists Y s.t.

$$\begin{aligned} \forall i \varphi_{e_i}(m_i, Y \oplus X), \\ \forall e, m (\varphi_e(m, Y \oplus X) \leftrightarrow (e, m) \in \{(e_i, m_i)\}_i). \end{aligned}$$

Thus $\{(e, m) : \varphi_e(m, Y \oplus X)\} \leq_T \{(e_i, m_i)\}_i \leq_T X'$. \square

We can make \mathcal{L}_2 -structures in a \mathcal{L}_2 -structure. These structures are called coded structure.

Definition(Coded Structure)

A coded structure is a pair $(N, \{W_n\}_n)$ s.t.

- (1) N has $0^N, +^N$ and so on,
- (2) each W_n is a subset of N .

We say a coded structure $(N, \{W_n\})$ is an ω -structure if $N = \mathbb{N}, 0^N = 0, +^N = +$ and so on.

Fact

There exists a Π_1^0 formula $\psi(X, M)$ s.t.

- (1) $\text{WKL}_0 \vdash \forall X \exists M \psi(X, M)$,
- (2) $\text{ACA}_0 \vdash \psi(X, M) \rightarrow M$ is a coded ω -model of WKL_0 including X .

Theorem (ACA_0)

For any $X \subseteq \mathbb{N}$, there exists a coded ω -model M of WKL_0 s.t. $X \in M$ and $\forall Y \in M (Y \oplus X)' \leq_T X'$.

[Proof] Let $X \subseteq \mathbb{N}$ and $\psi(X, M)$ be a Π_1^0 formula as in the previous fact. Then we can define an X -computable tree T s.t. $M \in [T] \leftrightarrow \psi(X, M)$. Since $\exists M \psi(X, M)$, $[T] \neq \emptyset$. By applying low basis theorem to T , we can get $M \in [T]$ s.t. $(M \oplus X)' \leq_T X'$. Now, if $Y \in M$ then $Y \leq_T M$ and hence $(Y \oplus X)' \leq X'$. \square

By similar argument, we can show the followings:

Theorem

ACA_0 proves relativized hyperimmune-free basis theorem:

$\forall X \forall T \leq_T X (|T| = \infty \rightarrow \exists Y \in [T])$

(Y is X -hyperimmune-free)).

Theorem (ACA_0)

For any $X \subseteq \mathbb{N}$, there exists a coded ω -model M of WKL_0 s.t. $X \in M$ and $\forall Y \in M (Y \text{ is } X\text{-hyperimmune-free})$.

Fact

Every noncomputable low set is not hyperimmune-free.

Corollary

Let M, M' be countable ω -models of WKL_0 s.t.

$X \in M \Rightarrow X$ is low , $X \in M' \Rightarrow X$ is hyperimmune-free.

Then $M \cap M' = \text{REC}$.

Corollary

In RCA_0 the following relations hold.

$$\text{WKL}_0 < \text{LBT} < \text{ACA}_0,$$
$$\text{WKL}_0 < \text{HFBT} < \text{ACA}_0.$$

Note: We have already shown that \leq holds.

[Proof] $M, M' : \omega$ -models of WKL_0 s.t.

M includes only low sets,

M' includes only hyperimmune-free sets.

(1st inequality) $M \models \text{WKL}_0 + \neg\text{HFBT}$,

$M' \models \text{WKL}_0 + \neg\text{LBT}$.

(2nd inequality) Since $\emptyset' \notin M, M'$, we have

$M \models \text{LBT} + \neg\Sigma_1^0\text{-CA}$,

$M' \models \text{HFBT} + \neg\Sigma_1^0\text{-CA}. \quad \square$

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